



# RAMA UNIVERSITY

[www.ramauniversity.ac.in](http://www.ramauniversity.ac.in)

## FACULTY OF ENGINEERING & TECHNOLOGY

BCS-501    Operating System

Lecturer-20

Manisha Verma

Assistant Professor

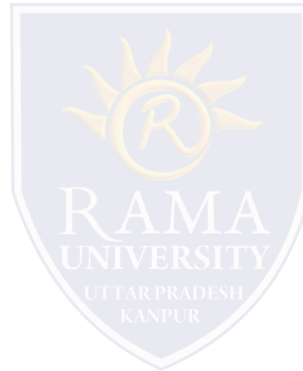
Computer Science & Engineering

**Contiguous Allocation**

**Hardware Support for Relocation and Limit Registers**

**Multiple-partition allocation**

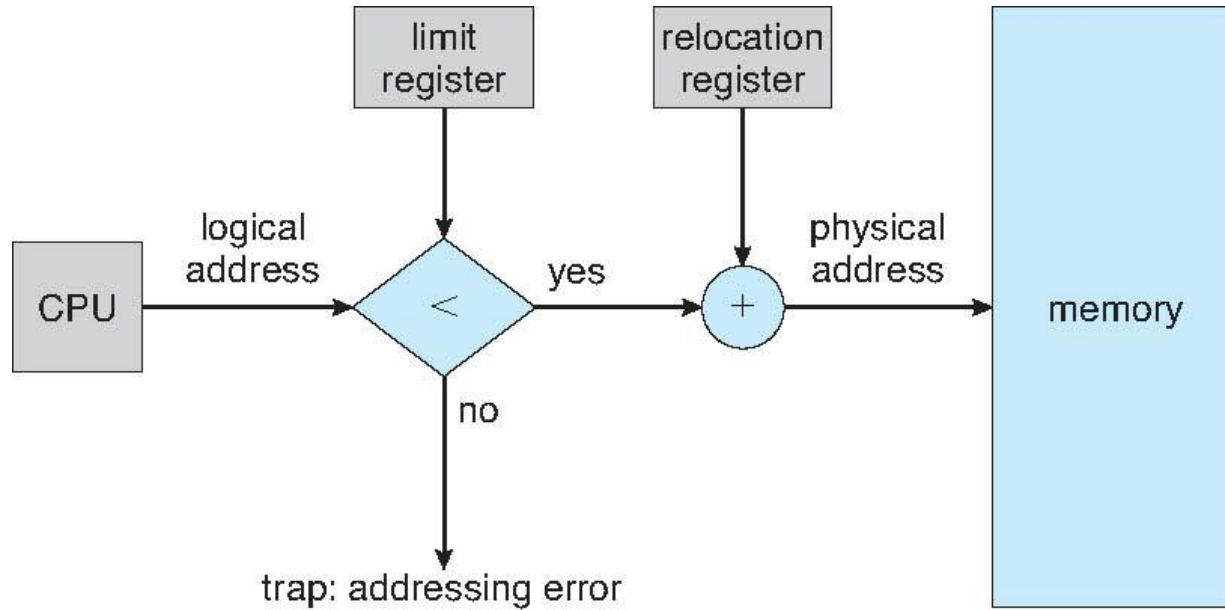
**Dynamic Storage-Allocation Problem**



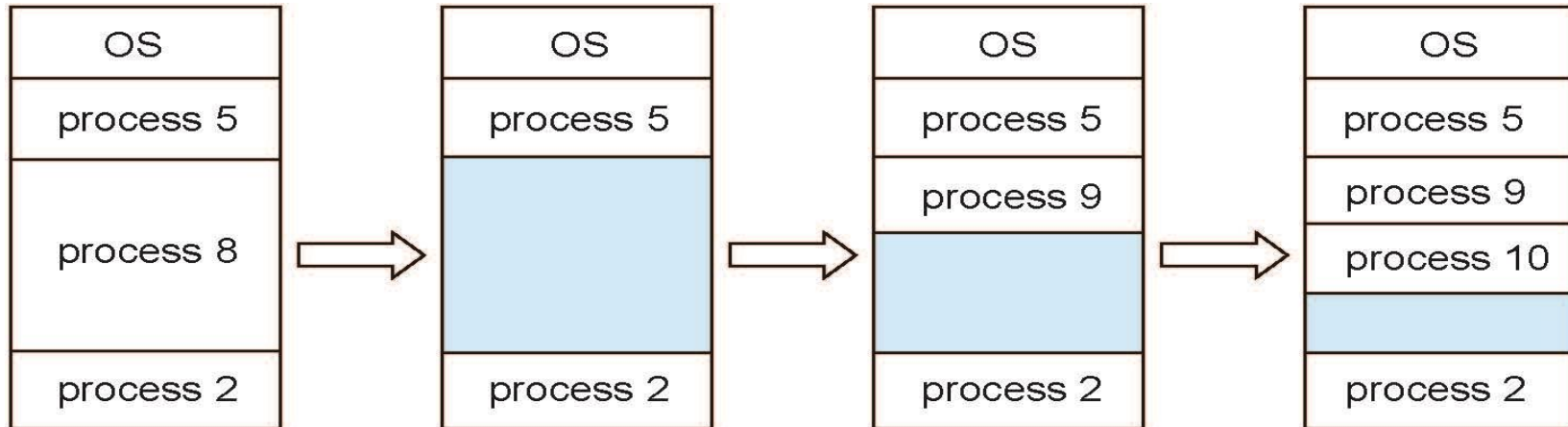
# Contiguous Allocation

- Main memory must support both OS and user processes
- Limited resource, must allocate efficiently
- Contiguous allocation is one early method
- Main memory usually into two partitions:
  - Resident operating system, usually held in low memory with interrupt vector
  - User processes then held in high memory
  - Each process contained in single contiguous section of memory
- Relocation registers used to protect user processes from each other, and from changing operating-system code and data
  - Base register contains value of smallest physical address
  - Limit register contains range of logical addresses – each logical address must be less than the limit register
  - MMU maps logical address dynamically
  - Can then allow actions such as kernel code being transient and kernel changing size

# Hardware Support for Relocation and Limit Registers



# Multiple-partition allocation



## •Multiple-partition allocation

- Degree of multiprogramming limited by number of partitions
- Variable-partition sizes for efficiency (sized to a given process' needs)
- Hole – block of available memory; holes of various size are scattered throughout memory
- When a process arrives, it is allocated memory from a hole large enough to accommodate it
- Process exiting frees its partition, adjacent free partitions combined
- Operating system maintains information about:
  - a) allocated partitions
  - b) free partitions (hole)

# Dynamic Storage-Allocation Problem

How to satisfy a request of size  $n$  from a list of free holes

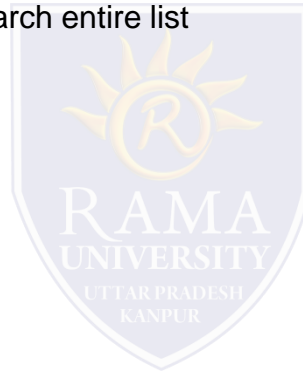
First-fit: Allocate the first hole that is big enough

Best-fit: Allocate the smallest hole that is big enough; must search entire list, unless ordered by size

Produces the smallest leftover hole

Worst-fit: Allocate the largest hole; must also search entire list

Produces the largest leftover hole



# MCQ

A process is thrashing if.....

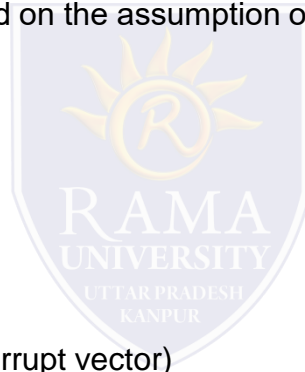
- A. it is spending more time paging than executing
- B. it is spending less time paging than executing
- C. page fault occurs
- D. swapping can not take place

Working set model for page replacement is based on the assumption of.....

- A. modularity
- B. locality
- C. globalization
- D. random access

The operating system is .....

- A. in the low memory
- B. in the high memory
- C. either a or b (depending on the location of interrupt vector)
- D. None of these



In contiguous memory allocation .....

- A. each process is contained in a single contiguous section of memory
- B. all processes are contained in a single contiguous section of memory
- C. the memory space is contiguous
- D. None of these

The relocation register helps in .....

- A. providing more address space to processes
- B. a different address space to processes
- C. to protect the address spaces of processes
- D. None of these

